Midterm

OCTOBER 20, 1992. 5:30-8:00 PM

There are seven problems for a total of 100 points. Good luck.

| Name: | | |
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| Name. | | |
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| Problem | Credit |
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| 1 | |
| 2 | |
| 3 | |
| 4 | |
| 5 | |
| 6 | |
| 7 | |
| Total | |

Are there more numbers than even numbers? To resolve this, either find a bijection between the naturals $\{0, 1, 2, \dots\}$ and the even, positive numbers $\{0, 2, 4, \dots\}$ or prove that none exists. What does this mean for the cardinality of the two sets: the set of naturals and the set of even naturals?

Consider the following function:

$$f(n) = \begin{cases} n/3 & \text{if } n \text{ is divisible by 3,} \\ \bot & \text{else.} \end{cases}$$

Give two different *total* functions g_1 and g_2 defined $N \to N$ which are extensions of f.

| Матн | 688: Theory of Computability and Complexity4 |
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| 3. (2 | 20 Points.) |
| | or each of the following mark "u" if the set is uncountable, "c" if it is buntable. |
| | The set of all real numbers. |
| | The set of all integers, both positive and negative. |
| | Fix a real $\epsilon > 0$, positive but no matter how small. The set of all reals in the interval $[1, 1 + \epsilon]$. |
| | The set of all subsets of the natural numbers. |

The set of all $\it finite$ subset of the natural numbers.

Prove: It is undecidable whether an arbitrary while-program halts on some input value. I.e., the following function is not computable:

$$f(i) = \begin{cases} 1, & \text{if } DOM(\varphi_i) \neq \emptyset; \\ 0, & \text{otherwise,} \end{cases}$$

where φ_i is the *i*-th computable unary function in our standard enumeration.

Write a while-program which computes the maximum of two numbers. I.e., the program moves to X1 the larger of the values in X1 and X2:

$$X1 := max(X1, X2)$$

Use no macros!

We define a new binary operation on functions, called *meet*. Given two partial functions f and g, the meet of f and g, $f \wedge g$, is the largest partial function h such that $h \leq f$ and $h \leq g$. That is, if h' is another partial function such that $h' \leq f$ and $h' \leq g$, then $h' \leq h$. In fact,

$$(f \wedge g)(i) = \begin{cases} f(i) \text{ or } g(i) \text{ when both are defined and equal,} \\ \bot \text{ elsewise.} \end{cases}$$

Show that if f and g are computable, then $f \wedge g$ is computable.

Use Church's Thesis to give an informal but rigorous argument that if a set of numbers cannot be effectively generated, then it is undecidable whether an arbitrary number is in the set.